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TITLE: Decoding method for trellis codes with large free distances

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INVENTOR-INFORMATION:

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CLAIMS:

What is claimed is:

1. A decoding method for a trellis code T of which the encoding can be implemented by first using an encoder of a binary trellis code $C_{sup}(1)$ to encode an r -bit message $u(t)$ into an intermediate binary m -tuple $v_{sup}(1)(t)$ at a t -th time unit which is sequentially converted into binary m -tuples $v_{sup}(2)(t), \dots, v_{sup}(L)(t)$ and a signal point $v_{sup}(L+1)(t)$ in a signal space Ω through L ($L \geq 2$) processors, comprising the decoding steps of:

(a) at a $(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L))$ -th time unit, a processor $P_{sup}(L)$ determining a set $T_{sub.M.sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$ which consists of $2_{sup.m}$ metrics for the $2_{sup.m}$ possible values of $v_{sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$, based on part of the set $\{y(t+i): i \leq \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L)\}$ and part of the set $\{v_{sup}(L)(t-j): j \geq \text{LAMBDA}\}$; feeding the metric set $T_{sub.M.sup}(L)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(L-1))$ into a processor $P_{sup}(L-1)$; wherein $y(t+i)$ is a received symbol to be decoded, $\text{LAMBDA}_{sup}(1), \dots, \text{LAMBDA}_{sup}(L)$ are nonnegative constants and LAMBDA is a truncation length to be used in decoding $C_{sup}(1)$;

(b) for $l=L-1, L-2, \dots, 1$, a processor $P_{sup}(l)$ determining a set $T_{sub.M.sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$ which consists of $2_{sup.m}$ metrics for the $2_{sup.m}$ possible values of $v_{sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$ based on part of $\{T_{sub.M.sup}(l+1)(t+i): i \leq \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l)\}$ and part of the set $\{v_{sup}(l)(t-j): j \geq \text{LAMBDA}\}$; feeding the metric set $T_{sub.M.sup}(l)(t + \text{LAMBDA}_{sup}(1) + \dots + \text{LAMBDA}_{sup}(l-1))$ into a processor $P_{sup}(l-1)$; wherein $\text{LAMBDA}_{sup}(1)$ is a nonnegative constant;

(c) a processor $P_{sup}(0)$ recovering the transmitted symbol $u(t - \text{LAMBDA} + 1)$ and $v_{sup}(1)(t - \text{LAMBDA} + 1)$ by applying the Viterbi decoding algorithm to the trellis of $C_{sup}(1)$ and using part of the set $\{T_{sub.M.sup}(1)(t-i): i \geq 0\}$; feeding back the recovered $v_{sup}(1)(t - \text{LAMBDA} + 1)$ to $P_{sup}(1)$;

(d) for $l=1, 2, \dots, L-1$, $P_{sup}(l)$ processing the set $\{v_{sup}(l)(t-j): j \geq \text{LAMBDA} - 1\}$ and determining $v_{sup}(l+1)(t - \text{LAMBDA} + 1)$ which is then fed back to $P_{sup}(l+1)$.

2. A decoding method as in claim 1, wherein the encoder of said single binary trellis

code $C_{sup}(1)$ is replaced by the encoders of a plurality of binary trellis codes which together convert $u(t)$ into $v_{sup}(1)(t) = (v_{sub.1.sup}(1)(t), v_{sub.2.sup}(1)(t), \dots, v_{sub.m.sup}(1)(t))$, which is then processed by L multilevel delay processors and L signal mappers; and wherein said Viterbi decoding algorithm for $C_{sup}(1)$ used in the processor $P_{sup}(0)$ is replaced by a plurality of Viterbi decoding algorithms for said plurality of binary trellis codes.

3. A decoding method as in claims 1 or 2, wherein the encoding of said trellis code T can be implemented by the steps of:

i) using the encoder of a single binary trellis code $C_{sup}(1)$ to encode an r -bit message $u(t)$ at the t -th time unit so as to obtain an m -bit output $v_{sup}(1)(t) = (v_{sub.1.sup}(1)(t), v_{sub.2.sup}(1)(t), \dots, v_{sub.m.sup}(1)(t))$;

ii) processing said output of the encoder through a multilevel delay processor $Q_{sup}(1)$ to obtain a processed output $s_{sup}(1)(t) = (s_{sub.1.sup}(1)(t), s_{sub.2.sup}(1)(t), \dots, s_{sub.m.sup}(1)(t))$ such that ##EQU11## wherein $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$ are nonnegative constants;

iii) feeding said output of $Q_{sup}(1)$ into a signal mapper $S_{sup}(1)$ to select a signal point $w_{sup}(1)(s_{sup}(1)(t)) = v_{sup}(2)(t) = (v_{sub.1.sup}(2)(t), v_{sub.2.sup}(2)(t), \dots, v_{sub.m.sup}(2)(t))$;

iv) for $l=2, 3, \dots, L$, feeding the output $v_{sup}(1)(t)$ into a multilevel delay processor $Q_{sup}(l)$ to obtain a processed output $s_{sup}(l)(t) = (s_{sub.1.sup}(l)(t), s_{sub.2.sup}(l)(t), \dots, s_{sub.m.sup}(l)(t))$, such that ##EQU12## wherein $\lambda_{sub.1.sup}(l), \lambda_{sub.2.sup}(l), \dots, \lambda_{sub.m.sup}(l)$ are nonnegative constants, and then feeding $s_{sup}(l)(t)$ into a signal mapper $S_{sup}(l)$ to select a signal point $w_{sup}(l)(s_{sup}(l)(t)) = v_{sup}(l+1)(t)$, whereby if $l=L$. A decoding method as in claim 3, wherein the parameter $\lambda_{sup}(l)$ of the processor $P_{sup}(l)$ is no less than $\Sigma_{i=1}^m \lambda_{sub.i.sup}(l)$ for $1 \leq l \leq L$.

5. A decoding method as in claim 3, wherein said constants $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$ are distinct values and/or at least one of the constants $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m.sup}(1)$ is not zero.

6. A decoding method as in claim 3, wherein said constants $\lambda_{sub.1.sup}(1), \lambda_{sub.2.sup}(1), \dots, \lambda_{sub.m-1.sup}(1)$ are all identical to a constant $\lambda_{sup}(1)$ and $\lambda_{sub.m.sup}(1) = 0$.

7. A decoding method as in claims 1 or 2, wherein said binary trellis code or a plurality of binary trellis codes is/are a linear and time-invariant code or a plurality of linear and time-invariant codes respectively, i.e., a binary convolutional code or a plurality of binary convolutional codes.

8. A decoding method as in claims 1 or 2, wherein the signal space Ω is a signal constellation; then said method can be used to decode trellis coded modulation.

9. A decoding method as in claims 1 or 2, wherein the signal space Ω is a collection of binary m -tuples; then said method can be used to decode binary trellis codes.

10. A decoding method as in claims 1 or 2, wherein the r -bit information, $u(t)$, represents $m/m_{sup}(1)$ binary $(r \cdot \text{multidot} \cdot m_{sup}(1) / m)$ -tuples, whereby $m_{sup}(1)$ is a positive integer.

11. A decoding method as in claims 1 or 2, wherein the binary m -tuple $v_{sup}(1)(t)$, $1 \leq l \leq L$, can be decomposed into $m/m_{sup}(1)$ binary $m_{sup}(1)$ -tuples, wherein $m_{sup}(1)$ is a positive integer; and $v_{sup}(L+1)(t)$ represents $m/m_{sup}(L+1)$ signal points in Ω which consists of $2^{(m \cdot \text{sp} \cdot (L+1) \cdot \text{sup})}$ signal points, wherein $m_{sup}(L+1)$ is a positive integer.